OpenMP (with Labs)

Philip Y. Nee Cornell CAC

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Based on materials developed by Kent Milfeld at TACC

Multithreading

- A programming model that allows multiple threads to exist in a single process, sharing the same memory space
- Thread A single sequence of programming instructions
- Why multithreading?
 - Speedup For an application that scales well, the speedup increases linearly with number of threads, theoretically

Goal

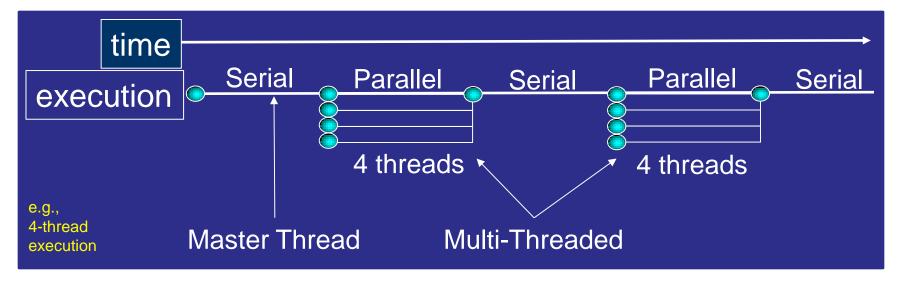
- In this talk, you will learn how to:
 - Parallelize portions of your application
 - Use the OpenMP API
 - Avoid common pitfalls in shared memory programming

What is OpenMP?

- OpenMP is an acronym for Open Multi-Processing
- An Application Programming Interface (API) for developing multithreading parallel program in shared memory architecture
- Three primary components of the API are:
 - Compiler Directives
 - Runtime Library Routines
 - Environment Variables
- OpenMP standard is maintained by the OpenMP Architecture Review Board
- http://www.openmp.org/ has the specification, examples, tutorials and documentation

OpenMP Fork-Join Parallelism

- Programs begin as a single process: master thread
- Master thread executes until a parallel region is encountered
 - Master thread creates (forks) a team of parallel threads
 - Threads in team simultaneously execute tasks in the parallel region
 - Team threads synchronize and terminate (join); master continues



Parallel Region: C/C++ and Fortran

LAB: OMP Hello World

```
#pragma omp parallel
    { code block
        a = work(...);
}
```

```
!$omp parallel
  code block
  call work(...)
!$omp end parallel
```

- Line 1 Team of threads is formed at parallel region
- Lines 2–3 Each thread executes code block and subroutine call, no branching into or out of a parallel region
- Line 4 All threads synchronize at end of parallel region (implied barrier)

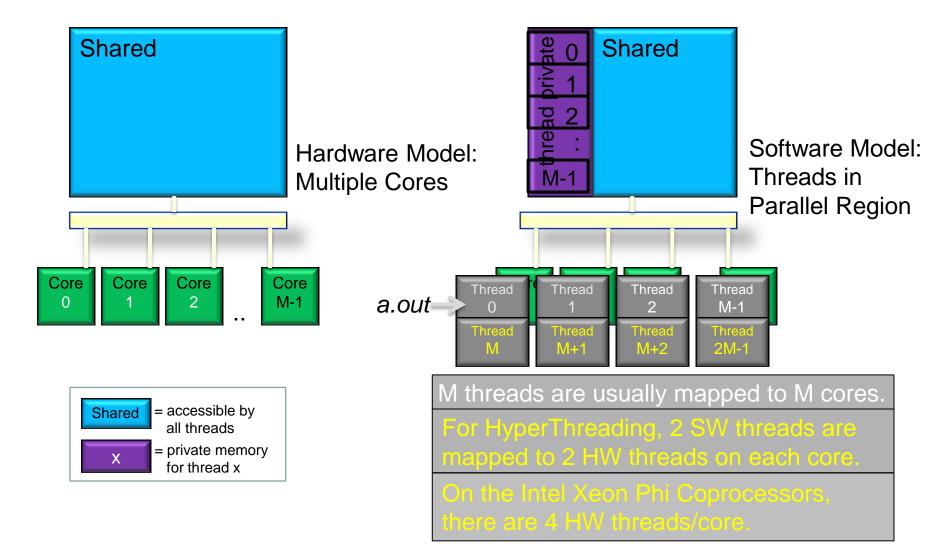
OpenMP: Multithreading

- All about executing concurrent work (tasks)
 - Tasks execute independently
 - Tasks access the same shared memory
 - Shared variable updates must be mutually exclusive
 - Synchronization through barriers

```
// repetitive work
#pragma omp parallel for
for (i=0; i<N; i++)
    a[i] = b[i] + c[i];</pre>
```

```
// repetitive updates
#pragma omp parallel for
for (i=0; i<N; i++)
    sum = sum + b[i]*c[i];</pre>
```

OpenMP: Shared Memory Systems



Thread Memory Access

- Every thread has access to "global" (shared) memory
 - All threads share the same address space
 - Threads don't communicate like MPI processes
- But need to avoid race conditions with shared memory. Examples:
 - 1. If multiple writers are going in no particular order, last writer "wins"
 - 2. A reader may either precede or follow a writer lack of synchronization
 - 3. Threads may overlap in a code block, causing conditions 1 and 2
- What do you with a race condition?
 - Don't introduce one in the first place: it's a bug, hard to debug
 - Impose order with barriers (explicit/implicit synchronization)
 - Use mutual exclusion (mutex) directives to protect critical sections,
 where one thread must run at a time (at a performance penalty)

Example of a Critical Section

Intended

Thread 0		Thread 1		Value
read	←			0
increment				0
write	\rightarrow			1
		read	←	1
		increment		1
		write	\rightarrow	2

Possible...

Thread 0		Thread 1		Value
				0
read	←			0
increment		read	←	0
write	\rightarrow	increment		1
		write	\rightarrow	1
				1

- In a critical section, need mutual exclusion to get intended result
- The following OpenMP directives prevent this race condition:

#pragma omp critical – for a code block (C/C++)

#pragma omp atomic – for single statements

OpenMP Directives

- OpenMP directives:
 - specify parallelism for shared-memory parallel (SMP) machines
 - Appear like comments in source code
- FORTRAN compiler: !\$OMP, C\$OMP, or *\$OMP use !\$OMP for free-format F90
- C/C++ compiler directive: #pragma omp

C/C++

```
#pragma omp parallel
    {...
}

#pragma omp parallel for
    for(...) {...
}
```

Fortran 90

```
!$OMP parallel
...
!$OMP end parallel
!$OMP parallel do
   DO ...
!$OMP end parallel do
```

Role of the Compiler

- OpenMP relies on the compiler to do the multithreading
 - Compiler recognizes OpenMP directives, builds in appropriate code
- A special flag is generally required to enable OpenMP
 - GNU: gcc -fopenmp
 - Intel: icc -openmp

OpenMP Compiler Directive Syntax

OpenMP Directives: Sentinel, construct, and clauses

```
#pragma omp construct [clause [[,]clause]...] C

!$omp construct [clause [[,]clause]...] F90
```

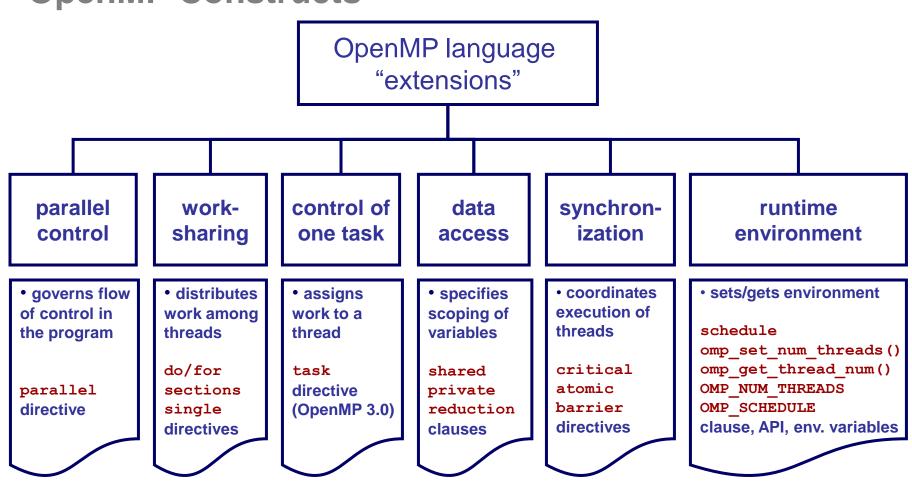
Example

```
#pragma omp parallel private(i) reduction(+:sum) C

!$omp parallel private(i) reduction(+:sum) F90
```

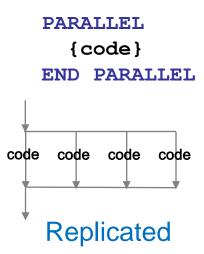
 Most OpenMP constructs apply to a "structured block", that is, a block of one or more statements with one point of entry at the top and one point of exit at the bottom.

OpenMP Constructs

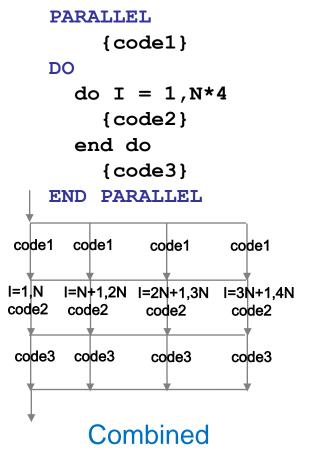


OpenMP Parallel Construct

- Replicated executed by all threads
- Worksharing divided among threads



```
PARALLEL DO
      do I = 1,N*4
         {code}
      end do
   END PARALLEL DO
     I=N+1,2N I=2N+1.3N I=3N+1.4N
code
     code
                     code
             code
      Worksharing
```



OpenMP Worksharing, Mutual Exclusion

Use OpenMP directives to specify worksharing in a parallel region, as well as mutual exclusion

parallel do/for
parallel sections

Directives can be combined, if a parallel region has just one worksharing construct.

Worksharing Loop: C/C++

```
1 #pragma omp parallel for
2 for (i=0; i<N; i++)
3 {
4    a[i] = b[i] + c[i];
5 }</pre>
```

General form:

```
#pragma omp parallel
{
    #pragma omp for
    for (i=0; i<N; i++)
    {a[i] = b[i] + c[i];}
}</pre>
```

Line 1 Team of threads formed (parallel region).

Lines 2–6 Loop iterations are split among threads. Implied barrier at end of block(s) {}.

Each loop iteration must be independent of other iterations.

Worksharing Loop: Fortran

General form:

```
!$omp parallel
!$omp do
do i=1,N
   a(i) = b(i) + c(i)
enddo
!$omp end parallel
```

Line 1 Team of threads formed (parallel region).

Lines 2–5 Loop iterations are split among threads.

Line 5 (Optional) end of parallel loop (implied barrier at enddo).

Each loop iteration must be independent of other iterations.

OpenMP Clauses

- Directives dictate what the OpenMP thread team will do
- Examples:
 - Parallel regions are marked by the parallel directive
 - Worksharing loops are marked by do, for directives (Fortran, C/C++)
- Clauses control the behavior of any particular OpenMP directive
- Examples:
 - 1. Scoping of variables: private, shared, default
 - 2. Initialization of variables: copyin, firstprivate
 - 3. Scheduling: static, dynamic, guided
 - 4. Conditional application: if
 - 5. Number of threads in team: num threads

LAB: Worksharing Loop

Private, Shared Clauses

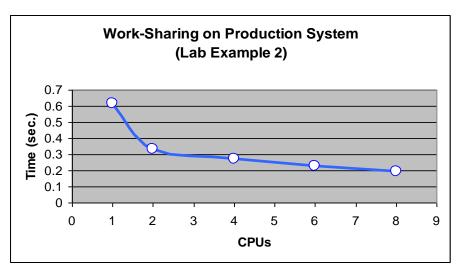
- Private: Variable is private to each thread
- Shared: Variable is shared among all threads
- In the following loop, each thread needs a private copy of temp
 - The result would be unpredictable if temp were shared, because each processor would be writing and reading to/from the same location

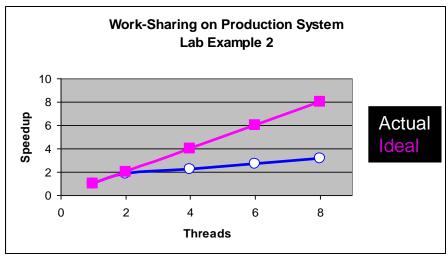
Worksharing Results

Speedup = cputime(1) / cputime(N)

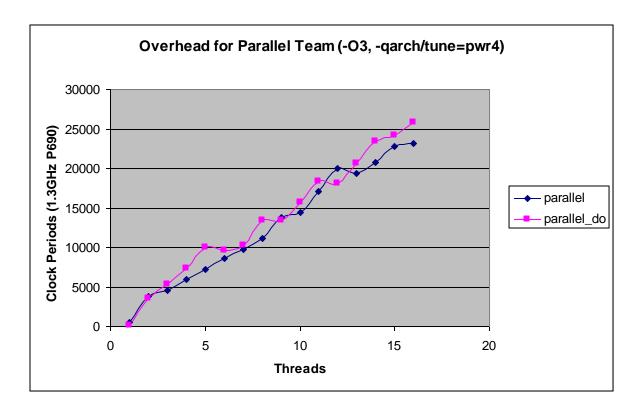
If work is completely parallel, scaling is linear.

Scheduling, memory contention and overhead can impact speedup and Mflop/s rate.





Overhead to Fork a Thread Team



Increases roughly linearly with number of threads

Merging Parallel Regions

Merging work-sharing constructs eliminates the overhead

```
!SOMP PARALLEL
!$OMP PARALLEL DO
                                         !$OMP DO
      do i=1,n
                                             do i=1,n
          a(i) = b(i) + c(i)
                                               a(i) = b(i) + c(i)
      enddo
                                             enddo
!$OMP END PARALLEL DO
                                         !$OMP END DO
!SOMP PARALLEL DO
                                         !$OMP DO
      do i=1,m
                                             do i=1,m
          x(i) = y(i) + z(i)
                                                 x(i) = y(i) + z(i)
      enddo
                                             enddo
!$OMP END PARALLEL DO
                                         !$OMP END DO
                                       !$OMP END PARALLEL
```

Runtime Library Functions

LAB: OMP Functions

<pre>omp_get_num_threads()</pre>	Number of threads in current team
<pre>omp_get_thread_num()</pre>	Thread ID, {0: N-1}
<pre>omp_get_max_threads()</pre>	Number of threads in environment, OMP_NUM_THREADS
<pre>omp_get_num_procs()</pre>	Number of machine CPUs
<pre>omp_in_parallel()</pre>	True if in parallel region & multiple threads executing
<pre>omp_set_num_threads(#)</pre>	Changes number of threads for parallel region, if dynamic threading is enabled

Environment Variables, More Functions

To control the OpenMP runtime environment

OMP_NUM_THREADS	Set to permitted number of threads: this is the value returned by omp_get_max_threads()
OMP_DYNAMIC	TRUE/FALSE for enable/disable dynamic threading (can also use the function below)

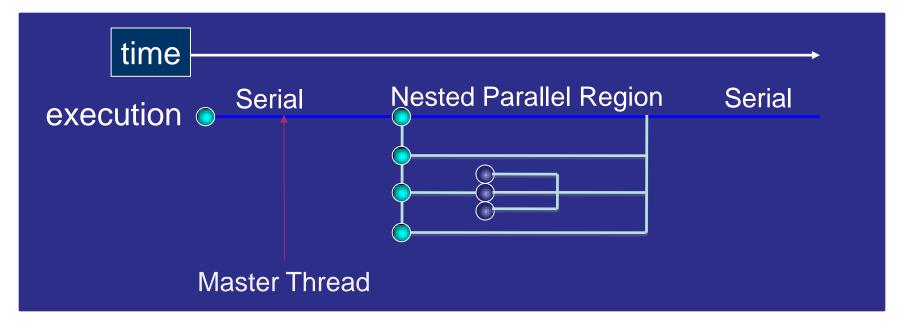
To enable dynamic thread count (not dynamic scheduling!)

<pre>omp_set_dynamic()</pre>	Set state of dynamic threading: if equal to "true", omp_set_num_threads() controls thread count
<pre>omp_get_dynamic()</pre>	True if dynamic threading is on

Additional OpenMP Features

- Schedule clause: specify how to divide work among threads
 schedule(static) schedule(dynamic, M)
- Reduction clause: perform collective operations on shared variables
 reduction (+:asum)
 reduction (*:aprod)
- Nowait clause: remove the barrier at the end of a parallel section
 for ... nowait
 end do nowait
- Lock routines: make mutual exclusion more lightweight and flexible
 omp init lock (var)
 omp set lock (var)

Loop Nesting



- OpenMP 3.0 supports nested parallelism, older implementations may ignore the nesting and serialize inner parallel regions.
- A nested parallel region can specify any number of threads to be used for the thread team, new id's are assigned.

MIC Programming with OpenMP

- Intel compiler, icc, is needed
- OpenMP pragma is preceded by MIC-specific pragma

```
- Fortran: !dir$ omp offload target(mic) <...>
- C: #pragma offload target(mic) <...>
```

- All data transfer is handled by the compiler
 - User control provided through optional keywords
- I/O can be done from within offloaded region
 - Data can "stream" through the MIC; no need to leave MIC to fetch new data
 - Also very helpful when debugging (print statements)
- Specific subroutines are automatically offloaded, including MKL subroutines

Example 1

2-D array (a) is filled with data on the coprocessor

Data management done automatically by compiler

- Memory is allocated on coprocessor for (a)
- Private variables
 (i,j,x) are created
- Result is copied back

```
use omp lib
                                  ! OpenMP
integer :: n = 1024
                                 1 Size
real, dimension(:,:), allocatable :: a ! Array
integer :: i, j
                                 ! Index
real
                              ! Scalar
            :: x
                            ! Allocation
allocate(a(n,n))
!dir$ omp offload target(mic) ! Offloading
!$omp parallel do shared(a,n), & ! Par. region
 private(x, i, j), schedule(dynamic)
do j=1, n
 do i=j, n
   x = real(i + j); a(i,j) = x
```

Example 2

Hand-coded routine, my sgemm:

specifies explicit in and out data
 movement

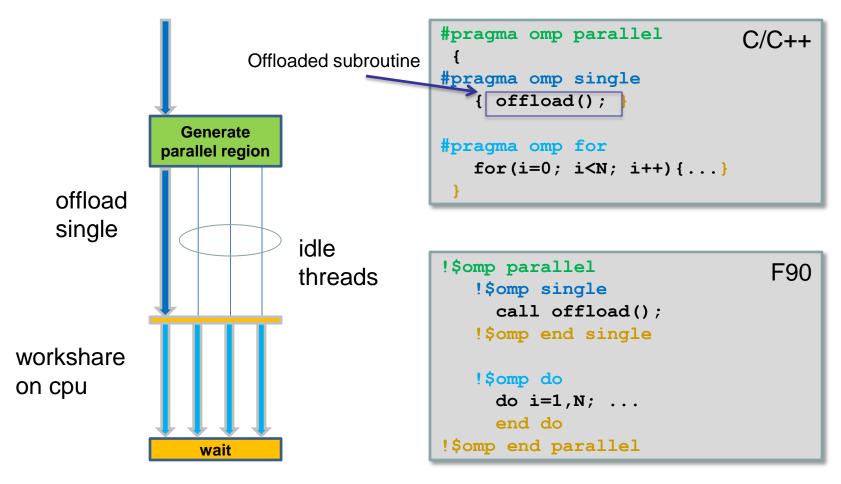
Use attributes to have routine compiled for the coprocessor, or link coprocessor-based MKL

LAB: Hand-coding vs. MKL

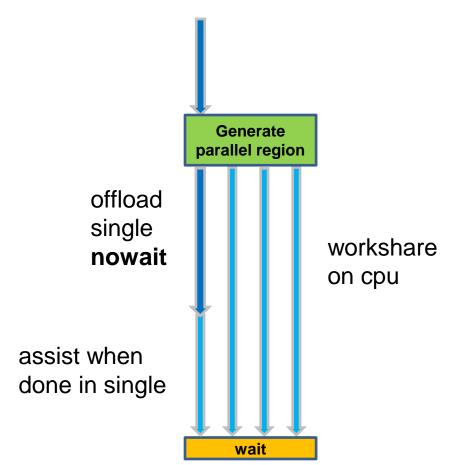
```
! snippet from the caller...
! offload MKL routine to accelerator
! offload hand-coded routine with data clauses
!dir$ offload target(mic) in(a,b) out(d)
call my_sgemm(d,a,b)
```

```
! snippet from the hand-coded subprogram...
!dir$ attributes offload:mic :: my_sgemm
subroutine my_sgemm(d,a,b)
real, dimension(:,:) :: a, b, d
!$omp parallel do
do j=1, n
    do i=1, n
    d(i,j) = 0.0
    do k=1, n
        d(i,j) = d(i,j)+a(i,k)*b(k,j)
    enddo; enddo; endo
end subroutine
```

Heterogeneous Threading, Sequential



Heterogeneous Threading, Concurrent



```
#pragma omp parallel
    {
    #pragma omp single nowait
        { offload(); }

#pragma omp for schedule(dynamic)
        for(i=0; i<N; i++){...}
}</pre>
```

```
!$omp parallel
   !$omp single
    call offload();
   !$omp end single nowait

   !$omp do schedule(dynamic)
    do i=1,N; ...
   end do
!$omp end parallel
```

Example 2

I/O from offloaded region:

- File is opened and closed by one thread (omp single)
- All threads take turns reading from the file (omp critical)

Threads may also read in parallel (not shown)

- Parallel file system
- Threads read parts from different targets

```
#pragma offload target(mic) //Offload region
#pragma omp parallel
  #pragma omp single /* Open File */
 printf("Opening file in offload region\n");
  f1 = fopen("/var/tmp/mydata/list.dat","r");
  #pragma omp for
  for(i=1;i<n;i++) {
    #pragma omp critical
    { fscanf(f1,"%f",&a[i]);}
    a[i] = sqrt(a[i]);
  #pragma omp single
 printf("Closing file in offload region\n");
  fclose (f1);
```